## **CLAYTON STATE** UNIVERSITY Design and testing of photon-based hardware random number generator

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# Random numbers & their uses

CYBERSECURITY

- Key generation
- Encryption algorithms

MACHINE LEARNING

- Model initialization
- Mutations
- Dropout selection

As well as physics simulations, lotteries, and video games.

# Pseudo-random number generators

#### Seed – an input that PRNGs use to define their randomness.

Seed for the world generator

Leave blank for a random seed

Figure 1. The seed-input box for the video game Minecraft.

Pros:

- Easily replicable results
- No external hardware necessary

Cons:

- Easily replicable results
- Not random enough for many use cases

# Photon-based solution to random number generation

Advantages of photon-based solutions

- Photon amplitude detection is an inherently chaotic process [1].
- The hardware involved is relatively cheap.



Figure 2. Concept design using an LED that pulses in sync with the reading of a photo-multiplier-tube.

#### Photon-based HRNG Prototype

Current prototype produces ~3500 amplitudes per second. Amplitudes are sent via serial and saved to a computer where they are processed.

Prototype hardware:

- Arduino Nano
- Multi-Pixel Photon Counter (MPPC) [2]
- Serially programmed power supply
- LED



Figure 3. Concept design using an LED that pulses in sync with the reading of a photo-multiplier-tube.

### Processing Amplitudes

### High/Low

if amp[i] > amp[i-1]: bitstream += 1
if amp[i] < amp[i-1]: bitstream += 0
else: skip</pre>

#### Even/Odd

if amp[i] % 2 == 0: bitstream += 1
else: bitstream += 0

#### Von Neumann's Procedure

- Search bitstream for pairs of 00 and 11.
- For each pair of 00 and 11, add a 0 to a new bitstream A and a corresponding 1 or 0 to bitstream a new B.
- For each non-matching pair, add a 0 to bitstream A and a 1 or 0 (depending on the first bit of the pair) to a new bitstream C.
- Recursively return the value of Neumann's procedure for each bitstream A, B, and C in the order of C + A + B.

#### Average & Standard Deviation Test

Using data collected over the course of 1 minute.

Pre-Neumann b	pitcount:	214676
Post-Neumann	bitcount:	208707

#### Pre-Neumann

This test is sensitive to data that has an inordinately unequal number of 1s and Os, although it does not determine whether the bitstream is necessarily random.

Average: 0.4987422907078574 Standard Deviation: 0.49999841816483426

#### Post-Neumann

Average: 0.5004000824121855 Standard Deviation: 0.49999983993403785

#### Monte Carlo Pi Estimation Test Results

Using data collected over the course of 1 minute.

Pre-Neumann bitcount: **214676** yielding **13417** coordinates Post-Neumann bitcount: **208707** yielding **13044** coordinates

Python also generates pi using its own random library but using the same number of coordinates as Pre-Neumann & Post-Neumann. This helps determine whether a poor-quality pi is due to poor quality data or too few coordinates.



**Figure 4.** Example visualization of the MCPE test

all\_dots += 1
if x^2 + y^2 <= 1: red\_dots += 1
pi = 4\*(red\_dots/all\_dots)</pre>

Figure 5. Psuedo-code describing how the MCPE test calculates pi

#### Pre-Neumann

Estimated Pi: 3.13572333607

#### Post-Neumann

Estimated Pi: 3.13002146581

#### Python's coords

Pre-neumman's bit total: 3.13453081911 Post-Neumann's bit total: 3.14872738424

#### The Fractional Line Symmetry Test

Count how frequently 1s or Os appear back-to-back for a determined length within the bitstream, called lines, where additional bits at the end of lines are added as fractional lines to the line total both horizontally and vertically for a stacked bitstream.

#### <u>Terms:</u>

Line: A single instance of back-to-back bits that meet the detect length.

**Detect Length:** How many times a bit must appear back-to-back before being counted to the line total.

Horizontal Search: Linearly search unmodified bitstream for lines.

**Vertical Search:** Stack bitstream into a square matrix and count number of lines top to down starting in the top right.

**Fractional Line:** Additional bits attached to an already full line that get added fractionally to the line total.

Example: Searching the bitstream 100010 horizontally for 0s with a detect length of 2 totalling 1.5 lines.



Figure 6. How the FLS test counts lines

#### **FLS Test Visualization**

Example Pre-Neumann bitstream visualization using High/Low processing.



Figure 7. Pre-Neumann High/Low
 bitstream visualization

#### **FLS Test Visualization**



**Figure 8.** Pre-Neumann High/Low bitstream horizontal FLS test.



Figure 9. Post-Neumann bitstream visualization vertical FLS test

### Vertical Line Counting

Original bitstream.



Add "nothing bits" to make bitstream stackable



Stack bitstream left to right, top to down into a square image and rotate Counter-clockwise





#### The FLS Test Results

Using data that was collected over the course of 1 minute.

Post-Neumann bitcount: 208707 Searching for: 0 Detect length: 4

Estimated Line Count:	8152.62
Horizontal Line Count:	8094.25
Vertical Line Count:	8190.5

#### Visualization



Figure 10. FLS test horizontal line visualization

Figure 11. FLS test vertical line visualization

The number of lines expected to be found can be calculated using  $\frac{n+1}{(2^{n+1})n} \times l$ 

where **n** is the detect length and **l** is the length of the bitstream.

### Conclusion

The photon-based solutions explored in this research have shown great potential as high speed solutions for high quality random number generation.

Future plans:

- Use a faster micro controller solution.
- Implement more tests for randomness such as the Ent test.
- Unify amplitude reading, processing, and testing into one program.
- House the HRNG.
- Secure the LED to the MPPC.
- Use voltage controller for the LED instead of PWM.



**Figure 12.** The current HRNG prototype alongside its future enclosure.

#### References

- [1] D. Beznosko et al., "Random Number Hardware Generator Using Geiger-Mode Avalanche Photon Detector", e -Print: 1501.05521, DOI: 10.22323/1.252.0049, PoS PhotoDet2015 (2016), 049
- [2] Hamamatsu Photonics K.K., 325-6, Sunayama-cho, Naka-ku, Hamamatsu City, Shizuoka Pref., 430-8587, Japan